# A robust measure of complexity

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# Roadmap

- 1 Research Question
- 2 Contribution
- 3 Theory
- Proof of concept



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# Roadmap

Research Question

# 2 Contribution

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### 5 Conclusion

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# Background

- Most Economists agree that (task) complexity is a key determinant of human behavior.
- It has the potential to explain mistakes that people systematically make
- At the same time, there is no consensus on what complexity is.
- Often, it is defined in a casual way

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# How is complexity formalized?

Direct approach: start with a definition, e.g.,

- Characteristics of lotteries (Huck & Weizsäcker, 1999; Fudenberg & Puri, 2023; Enke & Shubatt, 2023; Hu, 2023; de Clippel et al., 2025)
- Degree of contingent reasoning in mechanisms (Nagel & Saitto, 2025)
- How pronounced tradeoffs are (Shubatt & Yang, 2025)
- Through productivity of thinking about a task (Gabaix & Graeber, 2024)
- Signal-to-noise ratio (Goncalves, 2024)
- Revealed complexity: start with a measure/proxy, e.g.,

#### O Direct metrics:

- WTP to avoid a task (Oprea, 2020),
- response times (Wilcox, 1993; Goncalves, 2024),
- biometrics (van der Wel & van Steenbergen, 2018).
- Ø Behavioral metrics:
  - choice inconsistencies (Woodford, 2020).
- 8 Belief-based metrics:
  - expected accuracy (Agranov, Schotter & Trevino, 2025; Enke & Graeber, 2023; Enke, Graeber & Oprea, 2025; Hu, 2024; ...)

# Expected accuracy as a measure of complexity

• Expected accuracy: Probability to solve task correctly.

• Basic idea: Higher complexity := Lower expected accuracy

- Reasons to use it:
  - It is simple and intuitive!
  - ② Gaining momentum in the literature!
- On the flip side, there are two important caveats:
  - There are no choice theoretic foundations.
    - What are we actually measuring?
  - ② The induced complexity order depends on the size of the reward
    - Chances to solve task A are larger than task B, if reward is high
    - Chances to solve task A are smaller than task B, if reward is low
- Thus, the following question arises:

Is it a reasonable measure of complexity?

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# Expected accuracy depends on reward



• Reward affects attentions

(higher reward, more attention)

Attentions affects expected accuracy

(more attention, more likely to be correct)

• Hence, reward affects expected accuracy

# Expected accuracy depends on reward



• Reward affects attentions

(higher reward, more attention)

Attentions affects expected accuracy

(more attention, more likely to be correct)

• Hence, reward affects expected accuracy non-linearly

# Roadmap



# 2 Contribution

### B) Theory

### Proof of concept

### 5 Conclusion

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# Our approach

• We take a robust approach

 $\mathsf{Higher}\ \mathsf{complexity}\ :=\ \mathsf{Lower}\ \mathsf{expected}\ \mathsf{accuracy}\ \mathbf{for}\ \mathbf{every}\ \mathbf{reward}$ 

- Conservative (dominance) criterion: not all tasks will be ranked
- Complexity order:
  - (+) Intuitively appealing
  - (-) Incomplete

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# Preview of main conclusions

#### • Theoretical conclusions:

- Condition for any reasonable definition of complexity (based on the appeal of our criterion)
- ② Degree of uncertainty is an essential part of complexity (when is an exam deemed complex?)

#### • Practical conclusions:

- Recently popular measure of complexity is validated (using a lab experiment)
- Elicit expected accuracy for more rewards (using strategy method)

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# Our relation to the literature

- Literature justifies belief-based measures of complexity based on common sense (Agranov, Schotter & Trevino, 2025; Enke & Graeber, 2023; Enke, Graeber & Oprea, 2025; Hu, 2024; ...)
  - We provide theoretical foundations
- Literature takes complexity almost as a synonym to difficulty in the corresponding context (Oprea, 2024; Nagel & Saitto, 2025; Shubatt & Yang, 2025; Gabaix & Graeber, 2024; Goncalves, 2024; ...)
  - We identify degree of uncertainty as a novel channel
- Literature de facto postulates that complexity is a complete order (Oprea, 2020, 2024; Nagel & Saitto, 2025; Shubatt & Yang, 2025; Gabaix & Graeber, 2024; Goncalves, 2024; Woodford, 2020; Agranov, Schotter & Trevino, 2025; Enke & Graeber, 2023; Enke, Graeber & Oprea, 2025; Hu, 2024; ...)
  - No need to do so! We take a more conservative approach.

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# (Standard) theoretical framework

- Binary task  $S = \{s_0, s_1\}$
- Set of all tasks  $\mathcal{S}_0$
- Scores *Z* = {0, 1}
- Rewards  $X = [0, \infty)$
- Net utility  $v_{S}(x) := u_{S}(x, 1) u_{S}(0, 0) = \beta_{S}v(x)$ 
  - Task-specific subjective parameter (satisfaction)
  - Risk preferences are task-independent
  - Only for presentation purposes, fix  $eta_{\mathcal{S}}:=1$  for all  $\mathcal{S}\in\mathcal{S}_0$
- Prior belief  $\mu_{\mathcal{S}} \in [0,1]$  of  $s_1$ 
  - Novelty of our paper to let the prior vary across tasks
- Degree of uncertainty  $\eta_S = \frac{1}{\log 2} H(\mu_S)$ 
  - Task-specific subjective parameter (familiarity)
  - Consistent with information theory (Cover & Thomas, 2006)

# Degree of uncertainty



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# Degree of uncertainty



# Utility from answering correctly



# Utility from answering correctly



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# Attention

• Attention strategy: signal producing stochastic evidence



Each attention strategy is characterized by a (mean-preserving) distribution of posteriors: π ∈ Δ([0,1]) such that E<sub>π</sub>(q) = μ<sub>S</sub>.



- Attention has benefits and costs.
- Well-known that it is enough to focus on binary attention strategies (Matějka & McKay, 2015)

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# Expected benefit of attention graphically



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# Cost of attention

Posterior-separable cost of attention

$$\mathcal{C}_{\mathcal{S}}(\pi) = \kappa_{\mathcal{S}} \big( \mathbb{E}_{\pi}(c(q)) - c(\mu_{\mathcal{S}}) \big)$$

- Task-independent subjective parameter (cost of information processing)
- Task-specific objective parameter (difficulty)
- Solid theoretical foundations (Caplin et al., 2017; Tsakas, 2020; Zhong, 2022; Denti, 2022) and support by experimental findings (Dean & Neligh, 2024)
- Symmetry of *c* has been axiomatized (Hébert & Woodford, 2021) and particularly natural in binary tasks

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# Cost of information graphically



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- Agent's optimization problem  $\max_{\pi} (G_S(\pi) C_S(\pi))$
- Solved with concavification method (Aumann & Maschler, 1995; Kamenica & Gentzkow, 2011)



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# Optimal attention graphically



- Attention threshold  $q_S^{\times}$ :
  - decreasing in reward (x)
  - increasing in difficulty  $(\kappa_S)$

• Uncertainty threshold without reward  $\bar{\eta}_S = \frac{H(q_S^0)}{\log 2}$ 

# Attention map without reward



- Green area/large uncertainty ( $\eta_S > \bar{\eta}_S$ ): attention without reward
- Red area/small uncertainty  $(\eta_{S'} \leq ar{\eta}_{S'})$ : no attention without reward
- Without intrinsic incentives, the entire area is red

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# Expected accuracy



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# Robust definition of complexity

#### Definition

Task  $S \in \mathcal{S}_0$  is more complex than  $S' \in \mathcal{S}_0$  if

 $P(S,x) \leq P(S',x)$ 

for all  $x \ge 0$ . Then, we write  $S \succeq S'$ .

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# Trivial tasks

- Task  $S \in \mathcal{S}_0$  is trivial if  $S' \succeq S$  for all  $S' \in \mathcal{S}_0$ .
- The set of non-trivial tasks is denoted by  $\mathcal{S} \subseteq \mathcal{S}_0$ .



- The following are equivalent:
  - S is trivial
  - P(S, x) = 1 for all  $x \ge 0$
  - So easy that the state is learned with certainty (even without reward)

Characterization: Vector-valued representation (Ok, 2002)

#### Theorem (Identification)

For any pair  $S, S' \in S$ :

 $S' \succeq S \Leftrightarrow \phi(S') \ge \phi(S),$ 

where  $\phi_1(S) = \kappa_S$  and  $\phi_2(S) = \min\{\eta_S, \overline{\eta}_S\}$ .



• A task is complex when it is both difficult and unfamiliar.

(exam is complex when it is hard and not practiced).

• Difficulty remains the primary channel:  $\kappa_S > \kappa_{S'} \Rightarrow S' \not\geq S$ (even though not necessarily  $\kappa_S > \kappa_{S'} \Rightarrow S \succeq S'$ ).

• Without intrinsic incentives, we have  $\phi_2(S) = \eta_{S_{i,j}}$ 

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# Characterization of incompleteness

### Proposition (Single crossing)

Suppose that  $S, S' \in S$  are not  $\succeq$ -comparable, in that

- S more difficult than  $S': \phi_1(S) > \phi_1(S')$ ,
- S more familiar than  $S': \phi_2(S) < \phi_2(S')$ .

Then, there are two thresholds  $0 < x_1 < x_2 \le \infty$  such that:

(i) P(S, x) < P(S', x) for all  $x < x_1$ ,

- For small rewards, it is more likely to solve the difficult familiar task
- Not worthy paying attention, so the agent relies more on the prior
- (ii) P(S,x) > P(S',x) for all  $x_1 < x < x_2$ ,
  - For large rewards, it is more likely to solve the easy unfamiliar task
  - Worthy paying attention, so the agent relies more on the signal

(iii) P(S,x) = P(S',x) for all  $x \ge x_2$ .

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# Characterization of incompleteness

- S more difficult than  $S': \phi_1(S) > \phi_1(S')$
- S more familiar than  $S': \phi_2(S) < \phi_2(S')$



• Small rewards (red area): more likely to solve the difficult familiar task

- Large rewards (blue area): more likely to solve the easy unfamiliar task
- Very large rewards (green area): both tasks solved with certainty

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# Detour: Eliciting the complexity order

- Can we elicit the belief P(S, x)?
- Binarized scoring rule pays in probability units (Hossain & Okui, 2013):

Chances to win prizeWrong guess (z = 0)Correct guess (z = 1)Reported belief of  $z_1(R)$  $1 - \gamma R^2$  $1 - \gamma (1 - R)^2$ 

- Optimal report  $R(S,x) \neq P(S,x)$
- Identification problem due to state-dependent utilities (Tsakas, 2025)
  - She cares about the prize y and the outcome (x, z)
  - Even if we disregard hedging opportunities
- It doesn't matter: we actually care about  $\succeq$ , not about P(S, x):

# Proposition (Elicitation)

For every pair  $S, S' \in S$  and every  $x \ge 0$ :

 $P(S,x) \ge P(S',x) \iff R(S,x) \ge R(S',x).$ 

# Some practical considerations

• Robustness forces us to use strategy method:

- Elicit R(S, x) for multiple  $x \ge 0$
- Elicitation must take place before the task
- There are hedging opportunities
  - Usual problem (Blanco et al., 2010)
  - It can be solved by randomly paying for one  $x \ge 0$
- Alternative empirical strategy: Elicit belief about accuracy of others
  - Often simpler to implement
  - Similar idea in Bayesian markets (Baillon, 2017)
  - This is what we use in our experiment

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#### Proof of concept

- A panel with colored balls (red and blue) is drawn
  - Easy panel (100 balls): 51 of dominant color / 49 of other color
  - Difficult panel (400 balls): 201 of dominant color / 199 of other color
- This panel is drawn from a pool (all easy or all difficult)
  - Familiar task: 8 panels of one color / 2 panels of other color
  - Unfamiliar task: 5 panels of one color / 5 panels of other color
- Participants see the drawn panel and estimate the dominant color
- Two treatments: High reward (€10) / Low reward (€0.5)



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(H<sub>0</sub>) Sanity check: For all  $S \in \{EU, DU, EF, DF\}$ :  $P(S, \in 10) \ge P(S, \in 0.50)$ (H<sub>1</sub>) Basic hypothesis: For both  $x \in \{ \in 0.50, \in 10\}$  $P(EF, x) \ge P(EU, x) \ge P(DU, x)$  and  $P(EF, x) \ge P(DF, x) \ge P(DU, x)$ 

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Guess: how many participants were correct?



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 $\begin{array}{ll} (H_{2a}) & P(EU, \in 0.5) \geq P(DF, \in 0.5) \Rightarrow & P(EU, \in 10) \geq P(DF, \in 10) \\ (H_{2b}) & P(EU, \in 10) \leq P(DF, \in 10) \Rightarrow & P(EU, \in 0.5) \leq P(DF, \in 0.5) \end{array}$ 

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Guess: how many participants were correct?



### (H<sub>2</sub>) Additional hypotheses:

 $\begin{array}{l} (H_{2a}) \ P(EU, \in 0.5) \ge P(DF, \in 0.5) \ \Rightarrow \ P(EU, \in 10) \ge P(DF, \in 10) \\ (H_{2b}) \ P(EU, \in 10) \le P(DF, \in 10) \ \Rightarrow \ P(EU, \in 0.5) \le P(DF, \in 0.5) \end{array}$ 

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# Roadmap

- 1 Research Question
- 2 Contribution
- 3 Theory
- 4 Proof of concept



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# Overview of results

- Identification: Task A is more complex than task B if it is both more difficult and less familiar
  - Degree of uncertainty is a novel channel of complexity
  - Difficulty remains the primary channel
- Elicitation: Standard belief elicitation mechanisms reveal whether probability of solving A is larger or smaller than B, even though both might be misreported.
  - Not too difficult to elicit our measure
- **Validation**: Theoretical predictions corroborated in lab experiment.
- Completion (extra result): For non-comparable tasks A and B (viz., A is more difficult and more familiar than B), suppose that we start collecting data about B. Then, regardless where the data comes from, eventually A will certainly become more complex than B.
  - We do not need to know anything about the information source

# Thanks for your (in)attention ③

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